

STABILITY AND NUMERICAL ASPECTS OF LEAST SQUARES

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Monday, November 22, 2021

LOGISTICS

General announcements

- Assignment 6 to be posted... (grading traffic jam)
- 4 lectures left! (No lecture on Wednesday November 24, 2021)

Midterm 2

- 99% graded, grades released after Thanksgiving weekend

LOGISTICS

General announcements

- Assignment 6 to be posted... (grading traffic jam)
- 4 lectures left! (No lecture on Wednesday November 24, 2021)

Midterm 2

- 99% graded, grades released after Thanksgiving weekend
- Midterm solution during office hours on Tuesday November 23, 2021

Assignment 5

- Grading finalized

Assignment 6 and 7

- Posted this week, due date on Monday December 6 and Monday December 13

WHAT'S ON THE AGENDA FOR TODAY?

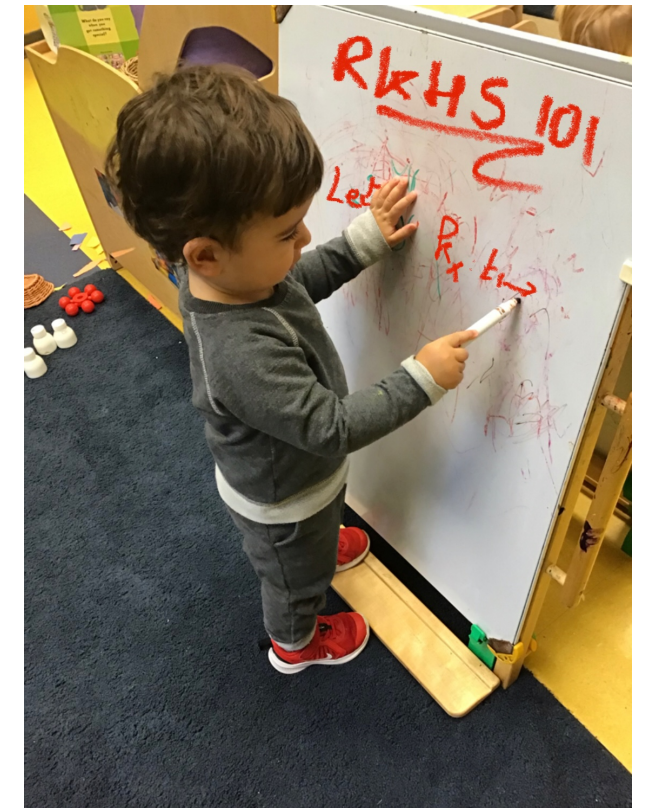
Last time:

- Stability of least squares

Today:

- More on stability and numerical considerations

Reading: lecture notes 14/15/16



Toddlers can do it!

STABILITY OF LEAST SQUARES

What if we observe $\mathbf{y} = \mathbf{A}\mathbf{x}_0 + \mathbf{e}$ and we apply the pseudo inverse? $\hat{\mathbf{x}} = \mathbf{A}^+\mathbf{y}$

We can separate the error analysis into two components

$$\hat{\mathbf{x}} - \mathbf{x}_0 = \underbrace{\mathbf{A}^+\mathbf{A}\mathbf{x}_0 - \mathbf{x}_0}_{\text{null space error}} + \underbrace{\mathbf{A}^+\mathbf{e}}_{\text{noise error}}$$

We will express the error in terms of the SVD $\mathbf{A} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^T$ With

- $\{\mathbf{v}_i\}_{i=1}^r$ orthobasis of $\text{row}(\mathbf{A})$, augmented by $\{\mathbf{v}_i\}_{i=1}^{r+1} \in \ker \mathbf{A}$ to form an orthobasis of \mathbb{R}^n
- $\{\mathbf{u}_i\}_{i=1}^r$ orthobasis of $\text{col}(\mathbf{A})$, augmented by $\{\mathbf{u}_i\}_{i=1}^{r+1} \in \ker \mathbf{A}^T$ to form an orthobasis of \mathbb{R}^m

The null space error is given by

$$\|\mathbf{A}^+\mathbf{A}\mathbf{x}_0 - \mathbf{x}_0\|_2^2 = \sum_{i=r+1}^n |\langle \mathbf{v}_i, \mathbf{x}_0 \rangle|^2$$

The noise error is given by

$$\|\mathbf{A}^+\mathbf{e}\|_2^2 = \sum_{i=1}^r \frac{1}{\sigma_i^2} |\langle \mathbf{e}, \mathbf{u}_i \rangle|^2$$

STABLE RECONSTRUCTION BY TRUNCATION

How do we mitigate the effect of small singular values in reconstruction?

$$\hat{\mathbf{x}} = \mathbf{V}\mathbf{\Sigma}^{-1}\mathbf{U}^T\mathbf{y} = \sum_{i=1}^r \frac{1}{\sigma_i} \langle \mathbf{y}, \mathbf{u}_i \rangle \mathbf{v}_i$$

Truncate the SVD to $r' < r$

$$\mathbf{A}_t \triangleq \sum_{i=1}^{r'} \sigma_i \mathbf{u}_i \mathbf{v}_i^T \quad \mathbf{A}_t^+ = \sum_{i=1}^{r'} \frac{1}{\sigma_i} \mathbf{u}_i \mathbf{v}_i^T$$

Reconstruct $\hat{\mathbf{x}}_t = \sum_{i=1}^{r'} \frac{1}{\sigma_i} \langle \mathbf{y}, \mathbf{u}_i \rangle \mathbf{v}_i = \mathbf{A}_t$

Error analysis:

$$\|\hat{\mathbf{x}}_t - \mathbf{x}_0\|_2^2 = \sum_{i=r+1}^n |\langle \mathbf{x}_0, \mathbf{v}_i \rangle|^2 + \underbrace{\sum_{i=r'+1}^r |\langle \mathbf{x}_0, \mathbf{v}_i \rangle|^2 + \sum_{i=1}^{r'} \frac{1}{\sigma_i^2} |\langle \mathbf{e}, \mathbf{u}_i \rangle|^2}_{\text{Tradeoff}}$$

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Error analysis:

$$\|\hat{\mathbf{x}}_t - \mathbf{x}_0\|_2^2 = \sum_{i=r+1}^n |\langle \mathbf{x}_0, \mathbf{v}_i \rangle|^2 + \sum_{i=r'+1}^r |\langle \mathbf{x}_0, \mathbf{v}_i \rangle|^2 + \sum_{i=1}^{r'} \frac{1}{\sigma_i^2} |\langle \mathbf{e}, \mathbf{u}_i \rangle|^2$$

STABLE RECONSTRUCTION BY REGULARIZATION

Regularization means changing the problem to solve

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \lambda \|\mathbf{x}\|_2^2 \quad \lambda > 0$$

The solution is

$$\hat{\mathbf{x}} = \underbrace{(\mathbf{A}^T \mathbf{A} + \lambda \mathbf{I})^{-1} \mathbf{A}^T \mathbf{y}}_{\textcircled{1}} = \mathbf{V}(\boldsymbol{\Sigma}^2 + \lambda \mathbf{I})^{-1} \boldsymbol{\Sigma} \mathbf{U}^T \mathbf{y}$$

(recall normal equations $(\mathbf{A}^T \mathbf{A} + \lambda \mathbf{I}) \mathbf{x} = \mathbf{A}^T \mathbf{y}$)

Scaling by $\frac{\sigma_i}{\sigma_i^2 + \lambda}$ for every $i \in \{1, \dots, r\}$

If $\sigma_i^2 \ll \lambda$ then $\frac{\sigma_i}{\sigma_i^2 + \lambda} \approx \frac{\sigma_i}{\lambda} \ll 1$

If $\sigma_i^2 \gg \lambda$ then $\frac{\sigma_i}{\sigma_i^2 + \lambda} \approx \frac{1}{\sigma_i} \ll 1$

Proof of $(A^T A + \lambda I) \tilde{A} y = V (\Sigma^2 + \lambda I)^{-1} \Sigma U^T y$

Recall that $\tilde{V} = [V \mid V_0] \in \mathbb{R}^{n \times n}$

Recall $A = U \Sigma V^T \in \mathbb{R}^{m \times n}$
 $A^T A \in \mathbb{R}^{n \times n}$

Introduce $\tilde{\Sigma} = \begin{pmatrix} \sigma_1 & & & & & \\ & \ddots & & & & \\ & & \sigma_r & & & \\ & & & & & 0 \\ 0 & & & & & \\ & & & & & \\ & & & & & \\ & & & & & 0 \end{pmatrix} \in \mathbb{R}^{n \times n}$

(in general $V^T V = I$ but $V V^T = I$)

Note that $\underline{A^T A + \lambda I}_n = \tilde{V} \tilde{\Sigma}^2 \tilde{V}^T + \lambda \tilde{V} \tilde{V}^T = \tilde{V} (\tilde{\Sigma}^2 + \lambda I) \tilde{V}^T$

Hence $(A^T A + \lambda I) \tilde{A} y = \tilde{V} \underbrace{(\tilde{\Sigma}^2 + \lambda I)^{-1}}_{(2)} \underbrace{\tilde{V}^T V}_{(1)} \Sigma U^T y$

Note $\tilde{V}^T V = [V \mid V_0]^T V = \begin{bmatrix} I_{r \times r} \\ \vdots \\ 0_{(n-r) \times r} \end{bmatrix}$ so that $\tilde{V} (\tilde{\Sigma}^2 + \lambda I)^{-1} \tilde{V}^T V = V (\Sigma^2 + \lambda I_r)^{-1}$

Error analysis:

$$\hat{x}_{\text{TK}} = \sum_{i=1}^r \frac{\sigma_i}{\sigma_i^2 + \lambda} \langle u_i, y \rangle v_i \quad \text{w/ } y = Ax_0 + e$$

$$= \sum_{i=1}^r \frac{\sigma_i}{\sigma_i^2 + \lambda} \langle u_i, Ax_0 \rangle v_i + \sum_{i=1}^r \frac{\sigma_i}{\sigma_i^2 + \lambda} \langle u_i, e \rangle v_i$$

$$\text{w/ } A = \sum_{i=1}^r \sigma_i u_i v_i^T$$

$$= \underbrace{\sum_{i=1}^r \frac{\sigma_i^2}{\sigma_i^2 + \lambda} \langle v_i, x_0 \rangle v_i}_{\textcircled{1}} + \underbrace{\sum_{i=1}^r \frac{\sigma_i}{\sigma_i^2 + \lambda} \langle u_i, e \rangle v_i}_{\textcircled{2}}$$

Recall $x_0 = \sum_{i=1}^n \langle v_i, x_0 \rangle v_i$; we will have the same kernel error as before

$\textcircled{1}$ is a source of "regularization" error w/ magnitude squared $\sum_{i=1}^r \left(1 - \frac{\sigma_i^2}{\sigma_i^2 + \lambda}\right)^2 |\langle v_i, x_0 \rangle|^2 = \sum_{i=1}^r |\langle v_i, x_0 \rangle|^2 \frac{\lambda^2}{(\sigma_i^2 + \lambda)^2}$

If $\sigma_i^2 \ll \lambda$ the error is $|\langle v_i, x_0 \rangle|^2$

If $\sigma_i^2 \gg \lambda$ the error is $\frac{\lambda^2}{\sigma_i^2} \ll 1$

$$\hookrightarrow \left\| \sum_{i=1}^r \langle v_i, x_0 \rangle v_i - \sum_{i=1}^r \frac{\sigma_i^2}{\sigma_i^2 + \lambda} \langle v_i, x_0 \rangle v_i \right\|_2^2$$

$\textcircled{2}$ is a source of noise error w/ magnitude squared $\sum_{i=1}^r \left(\frac{\sigma_i}{\sigma_i^2 + \lambda}\right)^2 |\langle u_i, e \rangle|^2 \leq \frac{1}{4\lambda} \sum_{i=1}^r |\langle u_i, e \rangle|^2 \leq \frac{1}{4\lambda} \|e\|_2^2$

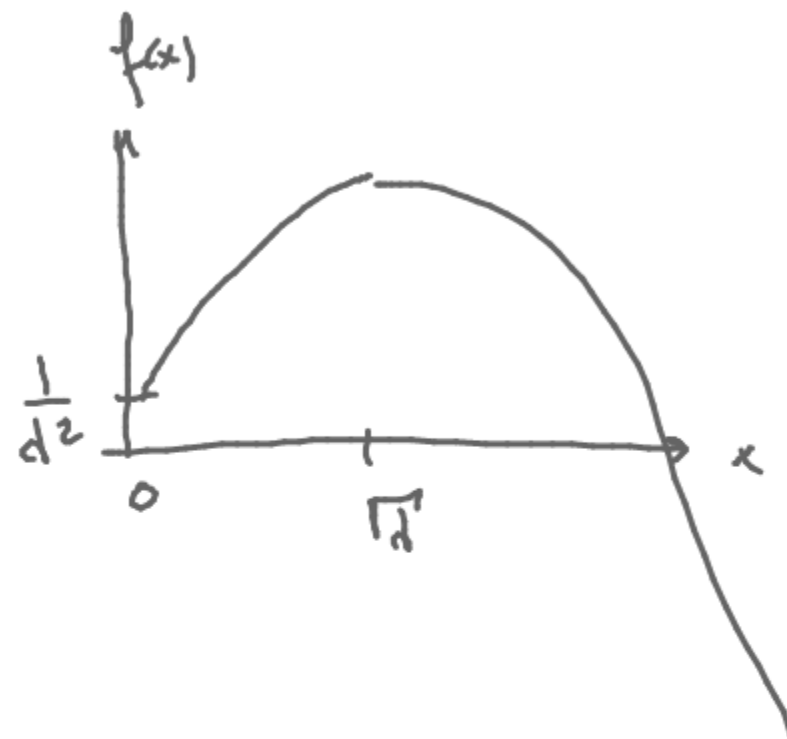
Proof of $\left(\frac{a}{a^2+b}\right)^2 \leq \frac{1}{4a}$

Set $f(x) = \left(\frac{x}{x^2+a}\right)^2 \quad x \in (0; +\infty) \quad f(0) = \frac{1}{a^2}$

$$f'(x) = 2 \frac{x}{x^2+a} \times \frac{x^2+a-2x^2}{(x^2+a)^2} = \frac{2x}{(x^2+a)^3} \times (a-x^2)$$

$$f(x) \geq 0 \Leftrightarrow a-x^2 \geq 0 \\ \Leftrightarrow x \leq \sqrt{a}$$

Note that $f(\sqrt{a}) = \left(\frac{\sqrt{a}}{2a}\right)^2 = \left(\frac{1}{2\sqrt{a}}\right)^2 = \frac{1}{4a}$



NUMERICAL METHODS

We have seen several solutions to systems of linear equations $\mathbf{Ax} = \mathbf{y}$ so far

- \mathbf{A} full column rank: $\hat{\mathbf{x}} = (\mathbf{A}^\top \mathbf{A})^{-1} \mathbf{A}^\top \mathbf{y}$
- \mathbf{A} full row rank: $\hat{\mathbf{x}} = \mathbf{A}^\top (\mathbf{A} \mathbf{A}^\top)^{-1} \mathbf{y}$
- Ridge regression: $\hat{\mathbf{x}} = (\mathbf{A}^\top \mathbf{A} + \delta \mathbf{I})^{-1} \mathbf{A}^\top \mathbf{y}$
- Kernel regression: $\hat{\mathbf{x}} = (\mathbf{K} + \delta \mathbf{I})^{-1} \mathbf{y}$
- Ridge regression in Hilbert space: $\hat{\mathbf{x}} = (\mathbf{A}^\top \mathbf{A} + \delta \mathbf{G})^{-1} \mathbf{A}^\top \mathbf{y}$

Extension: constrained least-squares

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{Ax}\|_2^2 \text{ s.t. } \mathbf{x} = \mathbf{B}\boldsymbol{\alpha} \text{ for some } \boldsymbol{\alpha}$$

- The solution is $\hat{\mathbf{x}} = \mathbf{B}(\mathbf{B}^\top \mathbf{A}^\top \mathbf{A} \mathbf{B})^{-1} \mathbf{B}^\top \mathbf{A}^\top \mathbf{y}$

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All these problems involve a symmetric positive definite system of equations.

- Many methods to achieve this based on matrix factorization

EASY SYSTEMS

Diagonal system

- $\mathbf{A} \in \mathbb{R}^{n \times n}$ invertible and diagonal
- $O(n)$ complexity

$$Ax = y = \begin{pmatrix} a_1 & & 0 \\ & \ddots & \\ 0 & & a_n \end{pmatrix} x = \begin{pmatrix} a_1 x_1 \\ \vdots \\ a_n x_n \end{pmatrix} = \begin{pmatrix} y_1 \\ \vdots \\ y_n \end{pmatrix} \quad \text{The solution is } x_i = \frac{y_i}{a_i}$$

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Orthogonal system

- $\mathbf{A} \in \mathbb{R}^{n \times n}$ invertible and orthogonal
- $O(n^2)$ complexity

$y = Ax$ so that $A^T y = x$ b/c $A^T A = I$ $O(n^2)$ is the cost of matrix multiplication

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Lower triangular system

- $\mathbf{A} \in \mathbb{R}^{n \times n}$ invertible and lower diagonal
- $O(n^2)$ complexity

$$Ax = y = \begin{pmatrix} a_{11} & & 0 \\ * & \ddots & \\ & & a_{nn} \end{pmatrix} \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} y_1 \\ \vdots \\ y_n \end{pmatrix}$$

ie.

$$\begin{array}{l} a_{11}x_1 = y_1 \quad \checkmark \\ a_{21}x_1 + a_{22}x_2 = y_2 \quad \checkmark \\ \vdots \\ a_{n1}x_1 + \dots + a_{nn}x_n = y_n \quad \checkmark \end{array} \quad \left. \vphantom{\begin{array}{l} a_{11}x_1 = y_1 \\ a_{21}x_1 + a_{22}x_2 = y_2 \\ \vdots \\ a_{n1}x_1 + \dots + a_{nn}x_n = y_n \end{array}} \right\} \text{substitution}$$

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Lower triangular system

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- $O(n^2)$ complexity

General strategy: factorize \mathbf{A} to recover some of the structures above

FACTORIZATIONS

LU factorization

$A \in \mathbb{R}^{n \times n}$ can be written $A = LU$ w/ $L \in \mathbb{R}^{n \times n}$ lower diagonal
 $U \in \mathbb{R}^{n \times n}$ upper diagonal

If A invertible then U and L are as well so that $Ax = b \Leftrightarrow L U x = b \Rightarrow x = U^{-1} L^{-1} b$ (*)

We can solve by $\left. \begin{array}{l} \textcircled{1} \text{ solve } Lw = b \text{ for } w \\ \textcircled{2} \text{ solve } Ux = w \text{ for } x \end{array} \right\} \text{triangular systems } O(n^2) \text{ (instead of } O(n^3) \text{ if inverting } A \text{)}$

FACTORIZATIONS

LU factorization

Cholesky factorization

QR decomposition

SVD and eigenvalue decompositions

COMPUTING EIGENVALUE DECOMPOSITIONS FOR SYMMETRIC MATRICES

Many techniques: we shall only discuss one based on *power iterations*